

# Strange Function

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Swiss Olympiad in Informatics

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```
// a >=0, b >= 0
long long funny(long long a, long long b) {
    long long x = 1;
    while (x < b) {
        x *= 2;
    }
    long long s = x % b;
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        a = (a & (x - 1)) + s * (a / x);
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# Easy subtasks

Task:

- Compute  $\text{funny}(a, b)$  for some specific  $(a, b)$  values.
- $\text{funny}(a, b)$  is guaranteed to return a result.



## Easy Subtasks: Obvious Solution

```
long long funny(long long a, long long b)
{ ... }
```

```
int main() {
    int T;
    cin >> T;
    for (int t = 1; t <= T; t++) {
        long long a, b;
        cin >> a >> b;
        cout << "Case_# " << t << " :_ ";
        cout << funny(a, b) << endl;
    }
}
```

# What's the point, then?

Full score only given for

- Constant time.
- Justification of correctness.

# Understanding the Code: x

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$x \leftarrow 2^i$ , for  $i$  such that  $x/2 < b \leq x$ .

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## Floor function

For integer  $n$ :

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Examples:

$$\lfloor 0.5 \rfloor = 0$$

$$\lfloor 1.6 \rfloor = 1$$

$$\lfloor 42.3 \rfloor = 42$$

$$\lfloor -1.3 \rfloor = -2$$



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## N.B. relation to C++

If  $a \geq 0$  and  $b > 0$ :

$a \bmod b = a \% b$ , where `%` is the corresponding operator in C++.

**Caution:** The condition is needed!

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$s \leftarrow x - b \cdot \underbrace{\lfloor x/b \rfloor}_{=1}$ .

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base 2

base 10:	0	1	2	3	4	5	6	7	8
base 2:	0	1	10	11	100	101	110	111	1000



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$a \& b$  keeps only the bits that are set in both  $a$  and  $b$ .

Example:

$$5 \& 6 = 101_{(2)} \& 110_{(2)} = 100_{(2)} = 4.$$

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- $x - 1$  in base 2:  $\dots 00 \underbrace{1 \dots 1}_{i \text{ times}}$ , therefore  $a \& (x - 1) = a \bmod x!$

# Understanding the Code

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**while**( $a \geq b$ ):  $a \leftarrow a - \underbrace{x \cdot \lfloor a/x \rfloor + x \cdot \lfloor a/x \rfloor}_{=0} - b \cdot \lfloor a/x \rfloor$ .

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- Upon termination,  $0 \leq a < b$ .
- `funny(a, b) == a%b!`

# Harder Subtasks

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- What about termination?  
 $x \leftarrow 2^i$ , for  $i$  such that  $x/2 < b \leq x$ .  
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**return**  $a$ .
- Nontermination for  $b \leq a < x$ .

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    if( $a' = a$ ) : return "NOTHING".  
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return  $a$ .
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Running time  $\mathcal{O}(\log(\min(a, b)))$ , not so nice to analyze.

# Optimizing running time

## Obvious Solution plus Trick

Run while loop until completion or until  $a$  no longer changes.

**if**( $a < b$ ): **return**  $a$ .

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We can show:  $b \leq a < x$  reached iff:  $a \bmod b + b < x$ .

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**if** ( $a < b$ ): **return**  $a$ .

$x \leftarrow 2^i$ , for  $i$  such that  $x/2 < b \leq x$ .

**if** ( $b = 0 \parallel a \bmod b + b < x$ ): **return** "NOTHING".

**return**  $a \bmod b$ .

Running time  $\mathcal{O}(\log(\log(\min(a, b))))$ .

(Binary search for exponent of  $x$  using bitshifts.)

Solution in  $O(1)$ Avoiding computation of  $x$ 

```
if( $b = 0 || a \geq b \&\&(b \&(b - 1)) \neq 0 \&\&((a \% b + b) \& \sim b) < b$ ) :  
    return "NOTHING".
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return  $a \bmod b$ .
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- More details in solution booklet.